Some Problems in Sanitizing Network Data

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Outline

- Problem statement and assumptions
- Network data sanitization
- Addresses and namespaces
 - IP address ranges and properties
 - Implications of finite namespaces
 - Nuances
- TCPsani
- Related work





Problem Statement

- collector collects data, gives it to analyst
- Some information in data is confidential
 - Sanitize it to remove enough information that sensitive data cannot be determined
- adversary tries to "reverse engineer" data to recover unsanitized information





Assumptions

- Adversary inserts markers into data as it is generated
- 2. Adversary has private information that, when combined with sanitized data, allows deduction of unsanitized data
- 3. Adversary can deduce unsanitized information directly from sanitized data





Our Assumptions

- 4. Adversary may have access to other sources that, when combined with sanitized data, allows deduction of unsanitized data
 - Privacy policy, threat model account for this
- 5. Collector makes data equally available to the analyst and adversary
 - Really, a worst case assumption





Why Sanitize?

- Need network traffic
 - Synthesize data
 - Capture, use reference data
 - Capture, sanitize data





Our Focus

- Packet headers
 - Care about MAC, IP, TCP/UDP headers
 - Overwrite application layer data
- Privacy policy: no IP address may be associated with an individual





Types of Anonymization

- Pseudo-anonymous
 - All instances of "John" become "Paul"
- Fully anonymous
 - John becomes "Paul", "George", "Ringo"
- Do mapping via tables, hash functions, etc.





Problems

- IP address ranges
 - Tables vs. hash functions
- Finite name spaces
 - Applying the pigeonhole principle
- Semantics
 - Sanitized data conveys contextual information not apparent from the data itself





IP Address Ranges

- Prevent mapping of 2 IP addresses to same target address
- Hash functions
 - Mutually distrusting collectors: risk dictionary attacks
- Tables
 - Mutually distrusting collectors can share portions of code book





Finite Namespaces

- Pseudo-anonymous mappings are permutations
- Fully anonymous mappings risk overflowing target namespace
 - Risks repetitions of sanitized names
 - "Bob" → any of 3 names
 - "Alice" → any of 10 names (more privacy)
 - Alternate approach: iterate sanitization





Semantics of Data

- Medical database
 - Not enough to sanitize names: other data might identify patient (date of birth, ZIP, gender)
 - Value of one field affects others; eg. ratio of height to weight significant
 - Names affect analysis; "Bob" unlikely to be at risk for pregnancy; "Shikigawa" more likely to be lactose-intolerant than "Smith"
- Analysis policy, metric must be explicit





The IP Version

- Packets to, from host all have port 53
 - Host is clearly a DNS server
 - Pseudo-anonymity preserves this
 - Full anonymity does not
- Pseudo-anonymity allows mapping the infrastructure





Moral

- In order to sanitize properly, three models required:
 - Threat model
 - Privacy policy
 - Analysis policy





TCPsani

- Inputs tcpdump savefile, sanitizes it, outputs in savefile format
 - Can be fed back to tcpdump, etc.
- Modified version of tcpdump that invokes Perl routines to sanitize
 - If you don't like ours, can write your own

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Supplied Versions

- Pseudo-anonymous mode:
 - Maps IP address byte by byte; IP address prefix determines map
 - Maps can be preconfigured to be shared
- Fully anonymous mode:
 - Map source region R to target region T
 - New address: pick empty slot in T, use it, add it to table





Related Work

- Classless IP address prefixes (tcpdpriv, CryptoPan)
 - Require trust among different collectors
- Network layer vs. application layer (Paxson, Pang)
- Pseudo-anonymous sanitization in file names and firewall logs (Sobirey, Fischer-Hübner, Rannenberg; Biskup, Flegel; Lundin, Jonsson)





Related Problems

- Privacy, analysis policies constrain inferences
- Database query-audit problem
- Cryptographic exchanges: crowds, anonymity set
- Semantics and structure (or lack thereof)
 - Preserving digital signatures of raw data to validate origin, integrity after sanitization





Conclusion

- Need explicit models to know what to do
- Related to several other, classic problems
- But critical we understand and make progress on it in order to further research and balance it with privacy needs
- Hard problem!





Concluding Thought

The truth about a man lies first and foremost in what he hides.

—Andre Malraux



