Lecture for January 22, 2016

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Overview

- Policies
- Trust
- Nature of Security Mechanisms
- Policy Expression Languages

Security Policy

- Policy partitions system states into:
 - Authorized (secure)
 - These are states the system can enter
 - Unauthorized (nonsecure)
 - If the system enters any of these states, it's a security violation
- Secure system
 - Starts in authorized state
 - Never enters unauthorized state

Confidentiality

- X set of entities, I information
- *I* has the *confidentiality* property with respect to *X* if no $x \in X$ can obtain information from *I*
- *I* can be disclosed to others
- Example:
 - *X* set of students
 - *I* final exam answer key
 - *I* is confidential with respect to *X* if students cannot obtain final exam answer key

Integrity

- X set of entities, I information
- *I* has the *integrity* property with respect to *X* if all $x \in X$ trust information in *I*
- Types of integrity:
 - Trust *I*, its conveyance and protection (data integrity)
 - *I* information about origin of something or an identity (origin integrity, authentication)
 - *I* resource: means resource functions as it should (assurance)

Availability

- X set of entities, I resource
- *I* has the *availability* property with respect to *X* if all *x* ∈ *X* can access *I*
- Types of availability:
 - Traditional: *x* gets access or not
 - Quality of service: promised a level of access (for example, a specific level of bandwidth) and not meet it, even though some access is achieved

Policy Models

- Abstract description of a policy or class of policies
- Focus on points of interest in policies
 - Security levels in multilevel security models
 - Separation of duty in Clark-Wilson model
 - Conflict of interest in Chinese Wall model

Mechanisms

- Entity or procedure that enforces some part of the security policy
 - Access controls (like bits to prevent someone from reading a homework file)
 - Disallowing people from bringing CDs and floppy disks into a computer facility to control what is placed on systems

Question

- Policy disallows cheating
 - Includes copying homework, with or without permission
- CS class has students do homework on computer
- Anne forgets to read-protect her homework file
- Bill copies it
- Who cheated?
 - Anne, Bill, or both?

Answer Part 1

- Bill cheated
 - Policy forbids copying homework assignment
 - Bill did it
 - System entered unauthorized state (Bill having a copy of Anne's assignment)
- If not explicit in computer security policy, certainly implicit
 - Not credible that a unit of the university allows something that the university as a whole forbids, unless the unit explicitly says so

Answer Part #2

- Anne didn't protect her homework
 Not required by security policy
- She didn't breach security
- If policy said students had to read-protect homework files, then Anne did breach security
 - She didn't do this

Types of Security Policies

- Military (governmental) security policy
 Policy primarily protecting confidentiality
- Commercial security policy
 Policy primarily protecting integrity
- Confidentiality policy

 Policy protecting only confidentiality
- Integrity policy
 - Policy protecting only integrity

Integrity and Transactions

- Begin in consistent state
 - "Consistent" defined by specification
- Perform series of actions (transaction)
 - Actions cannot be interrupted
 - If actions complete, system in consistent state
 - If actions do not complete, system reverts to beginning (consistent) state

Trust

Administrator installs patch

- 1. Trusts patch came from vendor, not tampered with in transit
- 2. Trusts vendor tested patch thoroughly
- 3. Trusts vendor's test environment corresponds to local environment
- 4. Trusts patch is installed correctly

Trust in Formal Verification

- Gives formal mathematical proof that given input *i*, program *P* produces output *o* as specified
- Suppose a security-related program *S* formally verified to work with operating system *O*
- What are the assumptions?

Trust in Formal Methods

- 1. Proof has no errors
 - Bugs in automated theorem provers
- 2. Preconditions hold in environment in which *S* is to be used
- 3. S transformed into executable S' whose actions follow source code
 - Compiler bugs, linker/loader/library problems
- 4. Hardware executes S' as intended
 - Hardware bugs (Pentium f00f bug, for example)

Types of Access Control

- Discretionary Access Control (DAC, IBAC)
 - Individual user sets access control mechanism to allow or deny access to an object
- Mandatory Access Control (MAC)
 - System mechanism controls access to object, and individual cannot alter that access
- Originator Controlled Access Control (ORCON)
 - Originator (creator) of information controls who can access information

Policy Languages

- Express security policies in a precise way
- High-level languages
 - Policy constraints expressed abstractly
- Low-level languages
 - Policy constraints expressed in terms of program options, input, or specific characteristics of entities on system

High-Level Policy Languages

- Constraints expressed independent of enforcement mechanism
- Constraints restrict entities, actions
- Constraints expressed unambiguously
 - Requires a precise language, usually a mathematical, logical, or programming-like language

Example: Ponder

- Security and management policy specification language
- Handles many types of policies
 - Authorization policies
 - Delegation policies
 - Information filtering policies
 - Obligation policies
 - Refrain policies

Entities

- Organized into hierarchical domains
- Network administrators
 - *Domain* is /NetAdmins
 - Subdomain for net admin trainees is
 - /NetAdmins/Trainees
- Routers in LAN
 - Domain is /localnet
 - Subdomain that is a testbed for routers is
 - /localnet/testbed/routers

Authorization Policies

• Allowed actions: netadmins can enable, disable, reconfigure, view configuration of routers

```
inst auth+ switchAdmin {
    subject /NetAdmins;
    target /localnetwork/routers;
    action enable(), disable(), reconfig(),
dumpconfig();
}
```

Authorization Policies

• Disallowed actions: trainees cannot test performance between 8AM and 5PM

```
inst auth- testOps {
    subject /NetEngineers/trainees;
    target /localnetwork/routers;
    action testperformance();
    when Time.between("0800", "1700");
}
```

Delegation Policies

• Delegated rights: net admins delegate to net engineers the right to enable, disable, reconfigure routers on the router testbed

inst deleg+ (switchAdmin) delegSwitchAdmin {
 grantee /NetEngineers;
 target /localnetwork/testNetwork/routers;
 action enable(), disable(), reconfig();
 valid Time.duration(8);
}

Information Filtering Policies

• Control information flow: net admins can dump everything from routers between 8PM and 5AM, and config info anytime

```
inst auth+ switchOpsFilter {
    subject /NetAdmins;
    target /localnetwork/routers;
    action dumpconfig(what)
        { in partial = "config"; }
        if (Time.between("2000", "0500")){
            in partial = "all"; }
}
```

Refrain Policies

• Like authorization denial policies, but enforced by the *subjects*: net engineers cannot send test results to net developers while testing in progress

```
inst refrain testSwitchOps {
    subject s=/NetEngineers;
    target /NetDevelopers;
    action sendTestResults();
    when s.teststate="in progress"
}
```

Obligation Policies

• Must take actions when events occur: on 3rd login failure, net security admins will disable account and log event

```
inst oblig loginFailure {
    on            loginfail(userid, 3);
    subject s=/NetAdmins/SecAdmins;
    target t=/NetAdmins/users ^ (userid);
    do            t.disable() -> s.log(userid);
}
```

Example

• Policy: separation of duty requires 2 different members of Accounting approve check

```
inst auth+ separationOfDuty {
    subject s=/Accountants;
    target t=checks;
    action approve(), issue();
    when s.id <> t.issuerid;
```

Low-Level Policy Languages

- Set of inputs or arguments to commands
 Check or set constraints on system
- Low level of abstraction
 - Need details of system, commands

Example: tripwire

- File scanner that reports changes to file system and file attributes
 - tw.config describes what may change /usr/mab/tripwire +gimnpsu012345678-a
 - Check everything but time of last access ("-a")
 - Database holds previous values of attributes

Example Database Record

/usr/mab/tripwire/README 0/. 100600 45763 1
917 10 33242 .gtPvf .gtPvY .gtPvY
0 .ZD4cc0Wr8i21ZKaI..LUOr3 .
0fwo5:hf4e4.8TAqd0V4ubv ?.... ...9b3
1M4GX01xbGIX00VuGo1h15z3 ?:Y9jfa04rdzM1q:eqt1AP
gHk ?.Eb9yo.2zkEh1XKovX1:d0wF0kfAvC ?
1M4GX01xbGIX2947jdyrior38h15z3 0

• file name, version, bitmask for attributes, mode, inode number, number of links, UID, GID, size, times of creation, last modification, last access, cryptographic checksums

Comments

- System administrators not expected to edit database to set attributes properly
- Checking for changes with tripwire is easy
 - Just run once to create the database, run again to check
- Checking for conformance to policy is harder
 - Need to either edit database file, or (better) set system up to conform to policy, then run tripwire to construct database