

Test and Set Solution

This algorithm solves the critical section problem for n processes using a Test and Set instruction (called TaS here). This instruction does the following function atomically:

```
function TaS(var Lock: boolean): boolean;
begin
    TaS := Lock;
    Lock := true;
end;
```

The solution is:

```
1 var waiting: shared array [0..n-1] of boolean;
2   Lock: shared boolean;
3   j: 0..n-1;
4   key: boolean;
5   ...
6 repeat (* process Pi *)
7   waiting[i] := true;
8   key := true;
9   while waiting[i] and key do
10    key := TaS(Lock);
11  waiting[i] := false;
12  (* critical section goes here *)
13  j := i + 1 mod n;
14  while (j <> i) and not waiting[j] do
15    j := j + 1 mod n;
16  if j = i then
17    Lock := false
18  else
19    waiting[j] := false;
20 until false;
```

lines 1-2: These are global to all processes, and are all initialized to false.

lines 3-4: These are local to each process i and are uninitialized.

lines 5-10: This is the entry section. Basically, $waiting[i]$ is true as long as process i is trying to get into its critical section; if any other process is in that section, then $Lock$ will also be true, and process i will loop in lines 8-9. Once process i can go on, it is no longer waiting for permission to enter, and sets $waiting[i]$ to false (line 10); it then proceeds into the critical section. Note that $Lock$ is set to true by the TaS instruction in line 9 that returns false.

lines 12-18: This is the exit section. When process i leaves the critical section, it must choose which other waiting process may enter next. It starts with the process with the next higher index (line 12). It checks each process to see if that process is waiting for access (lines 13-14); if no-one is, it simply releases the lock (by setting $Lock$ to false; lines 15-16). However, if some other process j is waiting for entry, process i simply changes $waiting[j]$ to false to allow process j to enter the critical section (lines 17-18).