Lecture 24 Outline

November 16, 2016

Reading: §16* Assignments: Homework 4, due Nov. 18; Lab 4, due Nov. 18

- 1. Access Control Lists
 - a. UNIX method
 - b. Full ACLs: describe, revocation issue
- 2. Capabilities
 - a. Capability-based addressing
 - b. Inheritance of C-Lists
 - c. Revocation: use of a global descriptor table
- 3. Lock and Key
 - a. Associate with each object a lock; associate with each process that has access to object a key (it's a cross between ACLs and C-Lists)
 - b. Example: cryptographic (Gifford). X object enciphered with key K. Associate an opener R with X. Then: **OR-Access**: K can be recovered with any D_i in a list of n deciphering transformations, so $R = (E_1(K), E_2(K), \dots, E_n(K))$ and any process with access to any of the D_i 's can access the file **AND-Access**: need all n deciphering functions to get K: $R = E_1(E_2(\dots E_n(K)\dots))$
 - c. Types and locks
- 4. Secret sharing
- 5. MULTICS ring mechanism
 - a. Rings, gates, ring-crossing faults
 - b. Used for both data and procedures; rights are REWA
 - (b_1,b_2) access bracket—can access freely; (b_3,b_4) call bracket—can call segment through gate; so if a's access bracket is (32, 35) and its call bracket is (36, 39), then assuming permission mode (REWA) allows access, a procedure in:
 - rings 0–31: can access a, but ring-crossing fault occurs
 - rings 32–35: can access a, no ring-crossing fault
 - rings 36–39: can access a, provided a valid gate is used as an entry point
 - rings 40-63: cannot access a
 - c. If the procedure is accessing a data segment *d*, no call bracket allowed; given the above, assuming permission mode (REWA) allows access, a procedure in:
 - rings 0-32: can access d
 - rings 33–35: can access d, but cannot write to it (W or A)
 - rings 36–63: cannot access d